# LINE-BLOCK GRAPHS AND CROSSING NUMBERS 

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The graph valued function namely the line-block graph $L_{b}(G)$ of a graph $G$ is the graph whose point set is the set of lines and blocks of $G$ and two points are adjacent if the corresponding blocks contain a common cutpoint of $G$ or one corresponds to a block $B$ of $G$ and other to a line $e$ of $G$ and $e$ is in $B$. In this paper, we establish a necessary and sufficient condition for graphs whose line-block graphs have crossing number $k, k=1,2,3$ or 4 .

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## Introduction

Throughout the paper, we only consider simple finite graphs. We follow [4, 5] for the terminology and notation. A block of a graph is a connected nontrivial graph having no cutpoint. The block graph $B(G)$ of a graph $G$ is the graph whose points are the blocks of $G$ and in which two points are adjacent whenever the corresponding blocks have a cutpoint in common. A graph $G$ is called a planar graph if $G$ can be drawn in a plane so that no two of its lines are cross each other. A graph that is not planar is called nonplanar. The crossing number $\operatorname{Cr}(G)$ of a graph $G$ is the minimum number of pairwise intersections of its lines when $G$ is drawn in the plane. Obviously $\operatorname{Cr}(G)=0$ if and only if $G$ is planar.

The line-block graph $L_{b}(G)$ of a graph $G$ is the graph whose point set is the set of lines and blocks of $G$ and two points are adjacent if the corresponding blocks contain a common cutpoint of $G$ or one corresponds to a block $B$ of $G$ and other to a line $e$ of $G$ and $e$ is in $B$. This concept was introduced by V. R. Kulli [6]. The crossing number of graph valued functions were studied in [1-3, 8-12].

The following theorem will be useful in the proof of our results.
Theorem 1.1[7] The line-block graph $L_{b}(G)$ of a graph $G$ is planar if and only if every cutpoint of $G$ is incident with at most 4 blocks.

## Lemmas

We start with preliminary lemmas, which are useful to prove our main theorems.
Lemma 2.1 If a graph $G$ has a cutpoint incident with six blocks, then $\operatorname{Cr}\left(L_{b}(G)\right) \geq 3$.

Proof. Suppose $G$ has a cutpoint $c$ incident with six blocks. Then the points corresponding to the blocks incident with cutpoint $c$ constitutes an induced subgraph $K_{6}$ in $L_{b}(G)$. It is known that $\operatorname{Cr}\left(K_{6}\right)=3$. Thus $\operatorname{Cr}\left(L_{b}(G)\right)=3$.

Suppose $G$ has a cutpoint incident with at most five blocks. Suppose $c_{1}$ is a cutpoint incident with five blocks. Then the points corresponding to the blocks incident with cutpoint $c_{1}$ constitutes an induced subgraph $K_{5}$ in $L_{b}(G)$. It is known that $\operatorname{Cr}\left(K_{5}\right)=1$. It follows that $\operatorname{Cr}\left(L_{b}(G)\right) \geq 3$.

Lemma 2.2 If a graph $G$ has two cutpoints each of which is incident with six blocks, then $\operatorname{Cr}\left(L_{b}(G)\right) \geq 6$.

Proof. Suppose $G$ has two cutpoints $c_{i}, i=1,2$ each of which is incident with six blocks. As in Lemma 2.1, points corresponding to six blocks incident with cutpoint $c_{i}$ forms an induced subgraph $K_{6}$ which has crossing number 3 . Since $G$ has two specified cutpoints $c_{1}$ and $c_{2}$. So $\operatorname{Cr}\left(L_{b}(G)\right)=6$.

Suppose $G$ has a cutpoint incident with at most five blocks. Suppose $c_{3}$ is a cutpoint incident with five blocks. Then the points corresponding to five blocks incident with cutpoint $c_{3}$ forms an induced subgraph $K_{5}$ in $L_{b}(G)$. It is known that $\operatorname{Cr}\left(K_{5}\right)=1$. It follows that $\operatorname{Cr}\left(L_{b}(G)\right) \geq 6$.

Lemma 2.3 If $G$ has $k(k \geq 1)$ cutpoints each of which is incident with five blocks and every other cutpoint of $G$ is incident with at most four blocks, then $\operatorname{Cr}\left(L_{b}(G)\right)=k$.

Proof. Suppose $G$ has $k(k \geq 1)$ cutpoints each of which is incident with five blocks. Then the points corresponding to five blocks incident with cutpoint $c_{i} ; 1 \leq i \leq k$ forms an induced subgraph isomorphic to $K_{5}$. It is known that $\operatorname{Cr}\left(K_{5}\right)=1$. Since $G$ has $k$ specified cutpoints $c_{i} ; 1 \leq i \leq k$, so $\operatorname{Cr}\left(L_{b}(G)\right)=k$ and a cutpoint of $G$ incident with at most four blocks forms $K_{2}$ or $K_{3}$ or $K_{4}$ as an induced subgraph, which is planar.

## Main results

Ingraphs with crossing number 1 or 2 .

Theorem 3.1. A graph $G$ has a line-block graph with crossing number $k ; k=1$ or 2 if and only if $G$ has exactly k cutpoints each of which is incident with five blocks and every other cutpoint of $G$ is incident with at most four blocks.

Proof. Suppose $\operatorname{Cr}\left(L_{b}(G)\right)=k ; k=1$ or 2 . Then $L_{b}(G)$ is nonplanar. By Theorem 1.1, $G$ has at least one cutpoint incident with at least five blocks. Assume $G$ has a cutpoint incident with six blocks. Then by Lemma 2.1, $\operatorname{Cr}\left(L_{b}(G)\right) \geq 3$, which is a contradiction. It implies that every cutpoint of $G$ is incident with at most five blocks and $G$ has at least one cutpoint incident with exactly five blocks. Suppose $c_{1}, c_{2}, \ldots, c_{r}$ be the cutpoints incident with five blocks in $G$. Then the points corresponding to five blocks incident with cutpoint $c_{i} ; 1 \leq i \leq r$ forms an induced subgraph isomorphic to $K_{5}$ in $L_{b}(G)$. Since $\operatorname{Cr}\left(K_{5}\right)=1$, and $L_{b}(G)$ has at least $r$ (line-disjoint) copies of $K_{5}, L_{b}(G)$ has at least $r$ crossings. If $r>k$, then $\operatorname{Cr}\left(L_{b}(G)\right) \geq k$, a contradiction. This proves that $r \leq k$. We consider two cases depending on the value of $k$. Assume $k=1$. Then $r=k$, for otherwise by Theorem 1.1, $\operatorname{Cr}\left(L_{b}(G)\right)=0$, a contradiction. Assume $k=2$. Suppose $r \leq 1$. Then $\operatorname{Cr}\left(L_{b}(G)\right) \neq k$, a contradiction. Therefore $r=k$. Thus $G$ has exactly $k$ cutpoints each of which is incident with five blocks and every other cutpoint
of $G$ is incident with at most four blocks, since a cutpoint of $G$ incident with at most four blocks forms $K_{2}$ or $K_{3}$ or $K_{4}$ as an induced subgraph, which is planar.

Conversely, suppose $G$ has exactly $k$ ( $k=1$ or 2 ) cutpoints of which is incident with five blocks and every other cutpoint of $G$ is incident with at most four blocks. Then by Lemma 2.3, $\operatorname{Cr}\left(L_{b}(G)\right)=k, k=1$ or 2.

We now characterize the line -block graphs with crossing number 3.
Theorem 3.2 A graph $G$ has a line-brock graph with crossing number 3 if and only if (i) or (ii) holds:
(i) $G$ has exactly three cutpoints each of which is incident with five blocks and every other cutpoint of $G$ is incident with at most four blocks.
(ii) $G$ has a unique cutpoint incident with six blocks and every other cutpoint of $G$ is incident with at most four blocks.

Proof. Suppose $\operatorname{Cr}\left(L_{b}(G)\right)=3$. Then clearly, $L_{b}(G)$ is nonplanar and by Theorem 1.1, $G$ has at least one cutpoint incident with at least five blocks. We consider two distinct cases:

Case 1. Assume $G$ has $k(k \neq 3)$ cutpoints each of which is incident with five blocks and every other cutpoint of $G$ is incident with at most four blocks. We consider two distinct subcases:

Subcase 1.1. Assume $G$ has $k(k \leq 2)$ cutpoints. Then by Theorem 3.1, $\operatorname{Cr}\left(L_{b}(G)\right) \leq 2$, which is a contradiction.

Subcase 1.2. Assume $G$ has $k(k \geq 4)$ cutpoints. Then by Lemma 2.3, $\operatorname{Cr}\left(L_{b}(G)\right) \geq 4$, again a contradiction. Thus from these two subcases we conclude that $G$ holds ( $i$ ).

Case 2. Assume $G$ has two cutpoints each of which is incident with six blocks. Then by Lemma 2.2, $\operatorname{Cr}\left(L_{b}(G)\right) \geq 6$, which is a contradiction. Thus $G$ holds (ii).

Conversely, suppose $G$ is a graph satisfying (i) or (ii). Then by Theorem 1.1, $L_{b}(G)$ has at least one crossing. We now show that its crossing number is 3 . First assume ( $i$ ) holds. Then by Lemma 2.3, $\operatorname{Cr}\left(L_{b}(G)\right)=3$. Now assume (ii) holds. Then by Lemma 2.1, $\operatorname{Cr}\left(L_{b}(G)\right) \geq 3$. But a cutpoint of $G$ incident with at most four blocks forms $K_{2}$ or $K_{3}$ or $K_{4}$ as an induced subgraph, which is planar. It follows that $\operatorname{Cr}\left(L_{b}(G)\right)=3$.

We now establish a necessary and sufficient condition for graphs whose line-block graphs having crossing number 4.

Theorem 3.3 A graph $G$ has a line-block graph with crossing number 4 if and only if (i) or (ii) holds:
(i) $G$ has exactly four cutpoints each of which is incident with five blocks and every other cutpoint of $G$ is incident with at most four blocks.
$G$ has a unique cutpoint incident with five blocks and six blocks respectively and every other cutpoint of $G$ is incident with at most four blocks.

Proof. Suppose $\operatorname{Cr}\left(L_{b}(G)\right)=4$. Then clearly, $L_{b}(G)$ is nonplanar and by Theorem 1.1, $G$ has at least one cutpoint incident with at least five blocks. We consider the following two cases :

Case 1. Assume $G$ has $k(k \neq 4)$ cutpoints each of which is incident with five blocks and every other cutpoint of $G$ is incident with at most four blocks. We consider the following two subcases:

Subcase 1.1. Assume $G$ has $k(k \leq 3)$ cutpoints. Then by Theorems 3.1 and 3.2, $\operatorname{Cr}\left(L_{b}(G)\right) \leq 3$, which is a contradiction.

Subcase 1.2. Assume $G$ has $k(k \geq 5)$ cutpoints. Then by Lemma 2.3, $\operatorname{Cr}\left(L_{b}(G)\right) \geq 5$, again a contradiction. Thus from these two subcases we conclude that $G$ holds ( $i$ ).

Case 2. In this case, we consider the following two subcases:
Subcase 2.1. Assume $G$ has two cutpoints each of which is incident with five blocks and a unique cutpoint incident with six blocks such that every other cutpoint of $G$ is incident with at most four blocks. Let $c_{1}$ and $c_{2}$ be the cutpoints incident with five blocks and let $c_{3}$ be the cutpoint incident with six blocks. Then by Theorem 3.1, the points corresponding to the blocks incident with cutpoints $c_{1}$ and $c_{2}$ constitute 2 crossings and by Lemma 2.1, the points corresponding to the blocks incident with cutpoint $c_{3}$ constitute at least 3 crossings in $L_{b}(G)$. Thus $L_{b}(G)$ has at least 5 crossings, which is a contradiction.

Subcase 2.2. Suppose $G$ has a unique cutpoint incident with five blocks and two cutpoints each of which is incident with six blocks such that every other cutpoint of $G$ is incident with at most four blocks. Let $c_{1}$ be the cutpoint incident with five blocks and let $c_{2}$ and $c_{3}$ be the cutpoints incident with six blocks. Then by Theorem 3.1, the points corresponding to the blocks incident with cutpoint $c_{1}$ forms 1 crossing and by Lemma 2.2, the points corresponding to the blocks incident with cutpoints $c_{2}$ and $c_{3}$ constitute at least 6 crossings in $L_{b}(G)$. Thus $L_{b}(G)$ has at least 7 crossings, again a contradiction. Thus from these two subcases we conclude that $G$ holds (ii).

Conversely, suppose $G$ is a graph satisfying $(i)$ or (ii), then by Theorem 3.2, $L_{b}(G)$ has more than 3 crossings. We now show that it has exactly 4 crossings. First assume (i) holds, then by Lemma 2.3, $\operatorname{Cr}\left(L_{b}(G)\right)=4$. Now assume (ii) holds. Let $c_{1}$ and $c_{2}$ be the cutpoints incident with five and six blocks respectively. By Theorem 3.1, the points corresponding to the blocks incident with cutpoint $c_{1}$ constitute one crossing in $L_{b}(G)$. Since $G$ has only one cutpoint incident with five blocks, By Lemma 2.1, the points corresponding to the blocks incident with cutpoint $c_{2}$ constitutes 3 crossing and a cutpoint of $G$ incident with at most four blocks forms $K_{2}$ or $K_{3}$ or $K_{4}$ as an induced subgraph, which is planar. Thus $L_{b}(G)$ has exactly 4 crossings.

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